

# PREPROCESSING IN SAT-BASED MULTI-OBJECTIVE OPTIMIZATION

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#### MOTIVATION

Maximum satisfiability (MaxSAT) solving has proven to be efficient for many real-world optimization problems

Many real-world applications have multiple conflicting objectives, calling for multi-objective (MO-)MaxSAT

**Preprocessing** (simplifying the instance before solving) is central in SAT, and becoming more popular in MaxSAT solving

#### WHY REDUNDANCY?



Instead of individual proofs  $( \dots )$ , uniform proofs of correctness and analysis of power  $(\rightarrow)$ 

#### **CONTRIBUTIONS**

- 3 distinct redundancy notions
- Lifting (Max)SAT preprocessing techniques to MO-MaxSAT
- Open-source preprocessor for MO-MaxSAT: MaxPre 2.1
- Empirical evaluation



#### **MULTI-OBJECTIVE MAXSAT**

MO-MAXSAT INSTANCE

- Constraints: clauses
- Min. objectives: pseudo-boolean  $O_i = \sum_x c_x^i \cdot x$

#### DOMINATING SOLUTIONS

 $\underline{\tau}$  weakly dominates  $\underline{\delta}$ : all objectives  $O_i(\underline{\tau}) \leq O_i(\underline{\delta})$ 

<u> $\tau$ </u> dominates  $\underline{\delta}$ : additionally one objective  $O_i(\underline{\tau}) < O_i(\underline{\delta})$ 

Every non-dominated solution is **Pareto-optimal** 



AIM Compute the non-dominated set (Pareto-optimal costs)  $\Rightarrow$  Preprocessing needs to preserve the **non-dominated** set

### - CONTRIBUTIONS

#### **REDUNDANCY IN MO-MAXSAT**

#### **Redundant clause** $C_{red}$ : does not change the non-dominated set.

For every solution  $\tau$  that does not satisfy  $C_{red}$ , there is a solution  $\delta$  that weakly dominates  $\tau$ and satisfies  $C_{\rm red}$ .

#### **Reconstructible clause** $C_{rec}$ : redundant clause satisfied by forcing a fixed set of literals $\omega$ .

The weakly dominating solution is always  $\delta = (\tau \setminus \neg \omega) \cup \omega$ .

#### **Literal-reconstructible clause**: reconstructible clause with a single literal in the set $\omega$ .

()

#### PARETO-MINIMAL CORRECTION SETS

**Pareto-MCS**: multi-objective extension of minimal correction set.

Pareto Sols. Pareto-MCSes Non-dom. Set



At least one Pareto-MCS per non-dominated cost tuple, but not all need to be preserved.

Literal-reconstructible clauses preserve Pareto-MCSes, but reconstructible clauses do not.

#### **RELATIONSHIPS BETWEEN REDUNDANCY NOTIONS**

A	All (	clauses			
	R	Redundant clauses			
		Reconstructible clauses			
		Literal-rec. clauses			

#### **EXAMPLE (REDUNDANT CLAUSES)**

 $F = \{(a_1 \lor a_2), (b_1 \lor b_2), (a_1 \lor b_1)\}$  $(a_1 \lor b_2), (a_2 \lor b_1), (a_2 \lor b_2), (a_3 \lor a_4) \}$  $O_1 = a_1 + a_2 + a_3 + a_4$  $O_2 = b_1 + b_2 + b_3$  $C_{\rm red} = (\neg a_2 \lor \neg b_2)$  $C_{\rm rec} = (\neg a_4), \quad \omega = \{a_3 \lor \neg a_4\}$  $C_{\text{l-rec}} = (\neg b_3), \quad I = \neg b_3$ 



#### MAXPRE 2.1 SUPPORTS LITERAL-RECONSTRUCTIBLE



- Blocked clause elimination
- Subsumption elimination
- Unit propagation\*
- Self-subsuming resolution
- Failed literal elimination\*
- Equivalent literal substitution\*
- ► TrimMaxSAT

\*on non-objective literals

RECONSTRUCTIBLE

## **EMPIRICAL EVALUATION**

**INSTANCE SIZE REDUCTION** 



#### PER-INSTANCE SOLVER PERFORMANCE IMPACT ON

Variables Clauses  $\sum$  Objective Coeffs. Pareto-MCSes

#### IMPACT ON SOLVER PERFORMANCE

	PackUP		LIDR		DAL				
# inst.	3692 (1420*)		366		96				
Solver	$\Delta \#$	$\Delta \sum t$	$\Delta \#$	$\Delta \sum t$ .	$\Delta \#$	$\Delta \sum t$			
BIOPTSAT (bi-objective optimization)									
LSU	+27	-13.7	+3	-18.3	—	_			
CG	+5	-6.4	$\pm 0$	-5.5	_	—			
Hybrid	+5	-13.6	$\pm 0$	-4.9	—	-			
Scuttle	+6	-40.4	-1	+7.2	+1	-0.6			
CLM									
CG	-5	+14.5	$\pm 0$	+1.4	-7	+4.7			
IHS	-19	-68.7	-19	-3.3	-7	-0.1			
LEXIMAXIST (leximax optimization)									

(Group-)subsumed label elimination CHANGE OBJECTIVES

- Unit propagation<sup>+</sup>
- Equivalent literal substitution<sup>+</sup>
- Intrinsic at-most-ones
- Binary core removal

<sup>+</sup>on objective literals



fraction remaining \*LEXIMAXIST (LSU) OBIOPTSAT (LSU) ×SCUTTLE



 $\Delta$  #: difference in number of solved instances

 $\Delta \sum t$ : difference in cumulative runtime over solved instances in 10<sup>3</sup>

seconds

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