



Certifying Pareto Optimality in Multi-Objective Maximum Satisfiability

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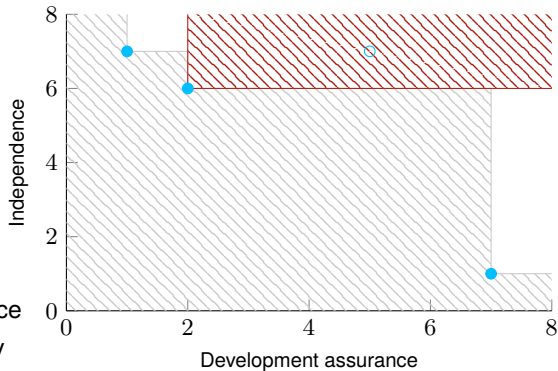


Multi-Objective Optimization

How to deal with conflicting objectives



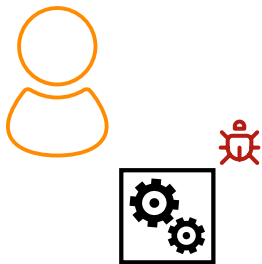
- ▶ Aviation: assurance level vs. independence
- ▶ Decision tree: accuracy vs. interpretability



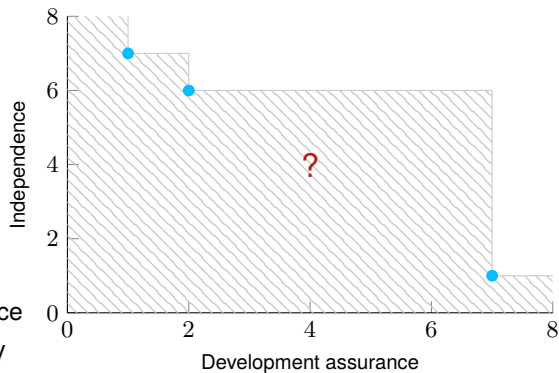


The Problem

Can you trust my solver?



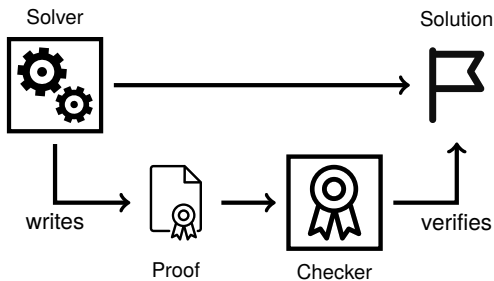
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Proof Logging

If we can't trust the solver, can we trust the result?



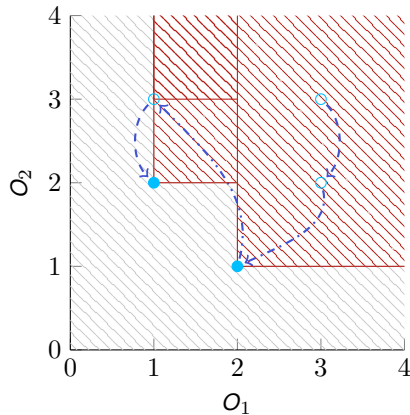
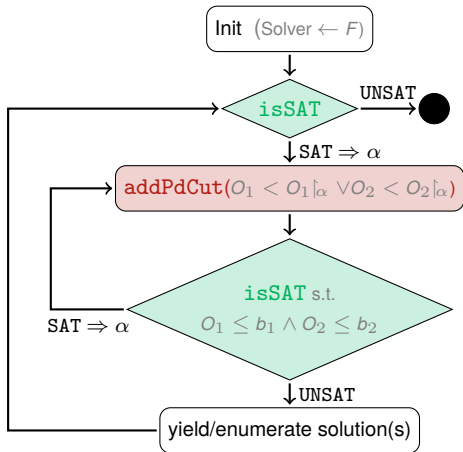
Paradigm	Proof format
SAT	DRAT
MaxSAT	VeriPB
MO-MaxSAT	This talk



The P -Minimal Algorithm

Multi-objective solution-improving search

[Soh et al. 2017]





The VERIPB Proof System

Pseudo-Boolean cutting planes proofs

[Bogaerts et al. 2023]

- ▶ Pseudo-Boolean Constraints
- ▶ Derive constraints by linear combination
- ▶ Redundant constraints
- ▶ Exclude solutions after finding them
- ▶ Single-objective optimization
- ▶ Preorder for expressing preference

minimize:

$$1x_1 + 2x_2 + 1x_3$$

prefer:

$$\bar{x}_4 \text{ over } x_4$$



$$[A] \quad 3x_1 + 2x_2 + 1x_3 \geq 3$$

$$[B] \quad \bar{x}_1 + x_4 \geq 1$$

$$[A+3B] \quad 2x_2 + 1x_3 + 3x_4 \geq 3$$

$$[\text{red}] \quad 1x_1 + 1x_2 + 2x_5 \geq 2$$

$$[\text{red}] \quad 1x_4 \geq 1$$

[solx] exclude x_1, x_2, x_3, x_4, x_5

[soli] better than $x_1, x_2, x_3, x_4, \bar{x}_5$



Multi-Objective Proof Setup

What VERIPB can already give us



prefer:
according to
Pareto optimality

reasoning steps from SAT solver
MO-specific reasoning

conclude as UNSAT

Guarantee

For each non-dominated point at least one solution explicitly appears in the proof

Syntactic restrictions

- ▶ First step in proof must load the Pareto order
- ▶ Order must never be changed



Encoding Pareto Dominance as a VERIPB Order

Telling VERIPB about the objectives

Given O_1, \dots, O_p

Required VERIPB order:

formula that is true iff α (weakly) dominates β

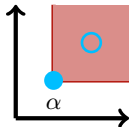
VERIPB Pareto order

$$O_i \upharpoonright_{\alpha} \leq O_i \upharpoonright_{\beta}, \quad \text{for } i = 1, \dots, p$$



Certifying Pareto Dominance Cuts

The building block for all algorithms



$$O_1 < O_1 \upharpoonright_{\alpha} \vee O_2 < O_2 \upharpoonright_{\alpha}$$

1. Reified objective CNF encoding
$$w_1 \Leftrightarrow O_1 < O_1 \upharpoonright_{\alpha}$$
2. Map each weakly dominated solution to α
(Redundant with α as witness)
3. Exclude α itself
4. Derive PD cut by combining previous two constraints

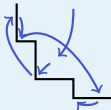


Proof Logging MO-MaxSAT Algorithms

Putting everything together

P-Minimal

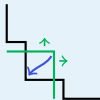
[Soh et al. 2017]



- ▶ SAT solver reasoning
- ▶ CNF objective encodings
- ▶ PD cuts

Lower-Bounding

[Cortes et al. 2023]



Upper-bounds irrelevant
→ same as *P*-Minimal

BIOPTSAT

[Jabs et al. 2024]

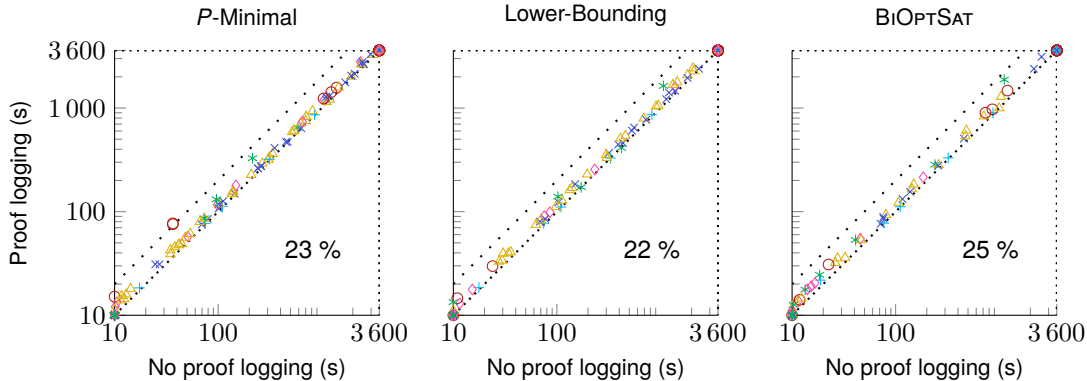


- ▶ Derive lower-bound on first objective
- ▶ Certify PD cut
- ▶ Strengthen PD cut based on known lower-bound



Proof Logging Overhead

How expensive is this





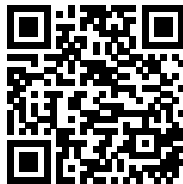
Proof Logging for Multi-Objective MaxSAT

Summary and conclusions

- ▶ MO-MaxSAT certifies that all non-dominated points were discovered
- ▶ Proofs in VERIPB format without extension
- ▶ Low overhead for proof logging
- ▶ Open-source implementation


Paper, slides, code, and contact:

`christophjabs.info/tacas25`





Bibliography I

-  [Bogaerts, Bart et al. \(2023\)](#). 'Certified Dominance and Symmetry Breaking for Combinatorial Optimisation.'. In: *JAIR* 77, pp. 1539–1589. doi: [10.1613/jair.1.14296](https://doi.org/10.1613/jair.1.14296).
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-  [Vandesande, Dieter et al. \(2022\)](#). 'QMaxSATpb: A Certified MaxSAT Solver'. In: *Logic Programming and Nonmonotonic Reasoning - 16th International Conference, LPNMR 2022, Genova, Italy, September 5-9, 2022, Proceedings*, pp. 429–442. doi: [10.1007/978-3-031-15707-3_33](https://doi.org/10.1007/978-3-031-15707-3_33). url: https://doi.org/10.1007/978-3-031-15707-3_33.

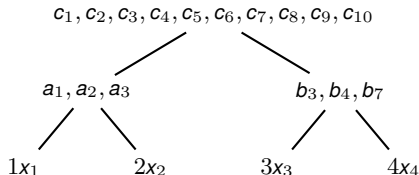


Certifying the Generalized Totalizer Encoding

[Vandesande et al. 2022]

Certified reified pseudo-Boolean CNF encodings

$$O = 1x_1 + 2x_2 + 3x_3 + 4x_4$$



1. Semantic definition of variables

$$a_2 \Leftrightarrow 1x_1 + 2x_2 \geq 2$$

(redundant because fresh variable)

2. Derive clauses from semantic definitions

$$(a_2^{\rightarrow} + b_3^{\rightarrow} + c_5^{\leftarrow} := (\overline{a_2} \vee \overline{b_3} \vee c_5))$$

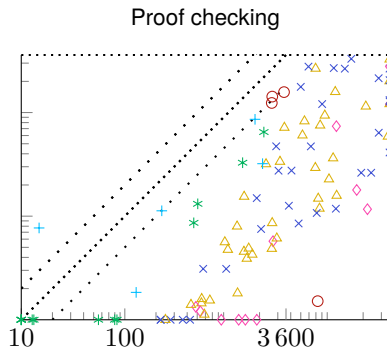
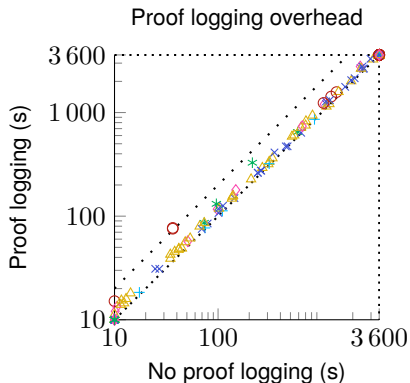
Our contribution

- ▶ This also works for the *generalized* totalizer
- ▶ Adapting process for complex core-boosted encodings



Proof Checking Performance

VERIPB still needs a lot of optimization work



Logging overhead:
23.3%

Checking:
47× longer